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Title: A Flexible Conservative Remapping Framework for Exascale Computing

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A Flexible Conservative Remapping Framework for Exascale Computing

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Now at Intuitive Surgical, Inc.

Introduction

- Remapping transfer of field data between computational meshes
- Remapping is used to transfer data between:
 - between two computational codes and their meshes,
 - between two physics modules and their meshes within a single code, or
 - from one mesh to another improved quality mesh as in Arbitrary-Lagrangian Eulerian (ALE) codes
- Desirable properties:
 - Conservative
 - Accurate
 - Bounds-preserving
 - Fast
 - Scalable



Conservative Remapping or Interpolation

- Compute an intensive quantity at entities of a target mesh from the source mesh such that a related extensive quantity is conserved across meshes
- For example, remap cell-centered values of density such that mass is conserved
- Often done by summing contributions of the extensive quantity from the source mesh to each target cell and converting back to an intensive quantity

Steps for Remapping Cell-based quantities

- Search: Identify the set of source mesh cell whose fields values potentially contribute to the target mesh cell
- Intersect: Find the moments of intersection (or volumes, centroids etc.) between the target entity and the source mesh entities that it overlaps
- Interpolate: Combine the field values from the overlapping source cells weighted using the moments of intersection to compute the target cell value

If the source cells overlapping a target cell are available on the same compute node, this is an embarassingly parallel algorithm



Portage Remapping Framework - Design

- Portage is a flexible framework for remapping
- Written using modern C++ (C++11 standard)
- Templated on almost all components allowing clients to assemble specialized remappers
- Design custom search, intersect and interpolate classes or use default ones
- Mesh and State class templating allows use with a wide set of applications
- Portage takes care of executing the algorithm with distributed and on-node parallelism



Portage Driver

Loosely speaking, the declaration of the driver is:



Portage Driver

Strictly speaking, the declaration is:

This is because the search, intersect and interpolate algorithms are also templated on dimension, the mesh wrappers and, where necessary, the state wrappers



Supplied Component Classes

- Mesh/State: Simple Mesh for cartesian meshes with scalar fields
- Optional interfaces to Jali (to be open-sourced) and FleCSI (open-source) mesh/state frameworks
- · Search: kd-tree search
- Intersect: R2D/R3D Open source intersection package (Devon Powell)
- Interpolation:
 - 1st-order accurate interpolation
 - 2nd-order accurate interpolation with Barth-Jespersen gradient limiting for bounds-preservation



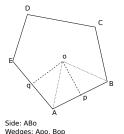
Handling Unstructured Polyhedral Meshes

- Mesh interface class assumes unstructured polyhedral meshes, possibly non-convex with curved faces
- Supplied search, intersect, interpolation designed for such meshes
- Basic functionality expected in mesh class:
 - Cells, face and node counts
 - cell o faces,dirs, cell o nodes, face o nodes, node o cells
 - cell ightarrow node-connected cells, node ightarrow cell-connected nodes
 - node coordinates, global IDs
 - parallel type of entities (owned, ghost)
 - boundary entities



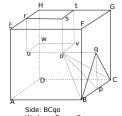
Handling Unstructured Polyhedral Meshes

- Handling non-convex cells requires:
 - decomposition of cells into tets (sides, wedges)
 - collecting wedges at a node to form the nodal control volume
- Portage provides this functionality, if client does not have it



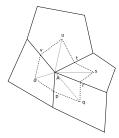
Polygon with subcell entities

Corner: pogA



Wedges: Bpqo, Cpqo Corner: uovwrstH

Polyhedron with subcell entities



Node with adjacent corners, wedges



Remapping of node centered fields

- Nodal quantities remapped as cell-centered quantities of dual mesh
- Dual control volumes formed by collecting the wedges around a node
- Works correctly for 1st order accurate remapping
- Incorrect for 2nd order accurate remapping because:
 - theory assumes linear reconstruction about the centroid of the control volume
 - node is not necessarily at the centroid of the dual control volume
- Also, nodal quantities are typically vector quantities like velocity
- Conservation of extensive quantity like momentum involves combination of node-centered (velocity) and cell-centered (density) quanities



Distributed memory parallelism

- Source and target mesh partitioning can be very different
- If source cells overlapping target cell are not on the same MPI rank, one has to fetch them
- Algorithm adapted from Parallel Rendezvous algorithm (Plimpton, et.al.) and Data Transfer Kit (Slattery, et.al.)
- Check overlap of bounding boxes of source (\mathcal{B}_s^i) and target (\mathcal{B}_t^l) partitions
- If $\mathcal{B}_t^i \cap \mathcal{B}_s^j \neq \emptyset$, request source partition j on rank i
- Send each source mesh partition and associated data to each rank that requested it - no communication needed subsequently
- Portage copies and transmits mesh/field data so that clients are not forced to redistribute
- Room for improvement send only the data that is necessary and avoid data copy if partitions are perfectly matched



On-node parallelism

- On-node parallelism achieved through NVidia Thrust parallel constructs like thrust::transform or thrust::for_each²
- Thrust can be directed to run the parallel constructs using an OpenMP, Intel TBB or CUDA backends
- Could also use Kokkos parallel constructs
- Abstracted out as Portage::transform so that we can call std::transform if Thrust not enabled
- Requires search, intersect and interpolate to have be written in functional style - no side effects
- Code is fully tested with OpenMP but only early version of 1st order remap tested with CUDA

²These constructs have also been accepted in the C++17 standard



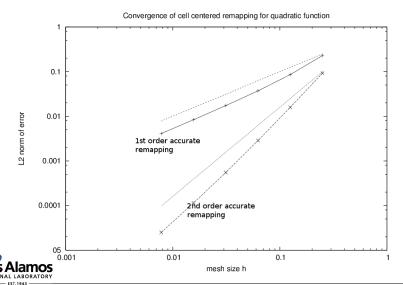
On-node parallelism - code outline

```
// Search for a list of overlap candidates for each target
Portage::transform(target_mesh_.begin(CELL), target_mesh_.end(CELL).
                   candidates.begin(), searchfunctor);
// Intersect each target cell with candidates cells to find
// intersection moments
Portage::transform(target_mesh_.begin(CELL), target_mesh_.end(CELL),
                   candidates, begin().
                   source_cells_and_weights.begin(),
                   intersectfunctor);
// Compute value of field on each target cell using the source moments
// and field values
Portage::transform(target_mesh_.begin(CELL), target_mesh_.end(CELL),
                   source_cells_and_weights.begin(),
                   target_field , interpolate);
```



Convergence of 1st and 2nd order accurate remaps

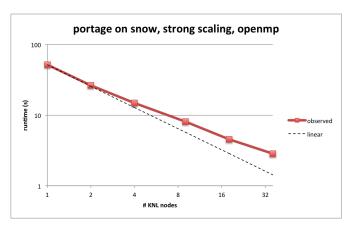
Quadratic analytical function



Results - OpenMP scaling

Intel Broadwell, 2 sockets/node, 18 cores/socket

Remap of single scalar field from a 36³ to a 48³ cartesian mesh.





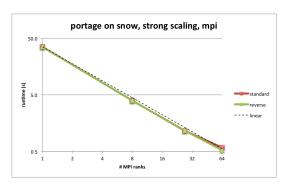
Results - MPI+OpenMP Strong Scaling

Intel Broadwell, 2 sockets/node, 1 MPI rank/socket, 18 cores/socket

Remap scalar field from 72³ to 96³ cartesian mesh

Standard partitioning refers to closely aligned partitioning for source and target mesh

Reverse partitioning reverses the MPI ranks for the source mesh to create a greater mismatch





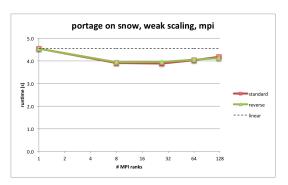
Results - MPI+OpenMP Weak Scaling

Intel Broadwell, 2 sockets/node, 1 MPI rank/socket, 18 cores/socket

Remap scalar field from 36³ to 48³ cartesian mesh *per MPI rank*

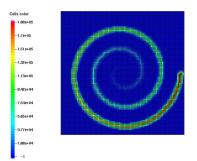
Standard partitioning refers to closely aligned partitioning for source and target mesh

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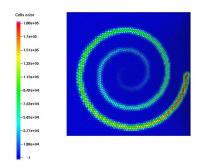




Results - 2D regular to polygonal remapping



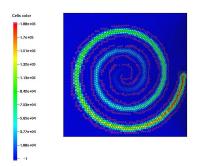
Source field



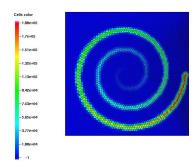
1st order remapping onto polygonal mesh



Results - 2D regular to polygonal remapping



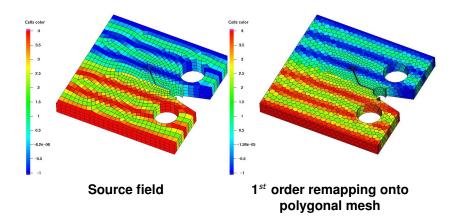
2nd order without limiter *Purple - out-of-bounds value*



2nd order with Barth-Jespersen limiting

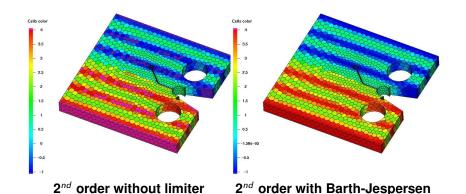


Results - 3D hexahedral to polyhedral mesh remapping





Results - 3D hexahedral to polyhedral mesh remapping



limiting



Purple - out-of-bounds values

On-going improvements

- Remapping of vector and tensor fields
 - Component-by-component simple
 - Preservation of invariants more complicated
- Improvement of intersection efficiency (most time spent here)
- "Smarter" gradient limiting at boundaries
- Minimizing data exchange between processors
- Third-order accurate reconstructions



Future work - Multi-material Remapping

- Remapping of quantities when souce cells may have multiple materials (hence multiple values of a variable)
- Only material values of field and volume fractions of materials in each cell known on source mesh
- Need to temporarily subdivide source cells into pure material polyhedra by interface reconstruction
- Material polyhedra can have general structure even for cartesian meshes
- Need to intersect target cells with material polyhedra not source cells
- Material-wise gradient (and higher order) reconstructions



Other future work

- Resurrect CUDA parallelism
- Mesh → particles → Mesh remapping
- Automatic data dependency resolution using FleCSI/Legion
- Automatic task parallelism through FleCSI/Legion
- Cartesian mesh, Spherical mesh and other specializations
- User Manual
- Wider deployment in codes



Availability

Open Source at http://www.github.com/laristra/portage
Mirrored periodically from internal LANL repository

